<https://github.com/roxanazachman01/FLCD>

**Documentation:**

**Statement:**

**Implement a parser algorithm (cont.)** - as assigned by the coordinating teacher, at the previous lab (PARSER LL(1))

**Deliverables**:

1. Algorithms corresponding to parsing table (if needed) and parsing strategy
2. Class ParserOutput - DS and operations corresponding to choice 2.a/2.b/2.c ([Lab 5](https://moodle.cs.ubbcluj.ro/mod/assign/view.php?id=2841)) (required operations: transform parsing tree into representation; print DS to screen and to file)

**Remark**: If the table contains conflicts, you will be helped to solve them. It is important to print a message containing row (symbol in LL(1), respectively state in LR(0)) and column (symbol) where the conflict appears. For LL(1), values (α,i) might also help.

**Implementation**:

**Parsing table algorithm:**

For each nonterminal we get the productions, and for each production we get the first set. For each terminal in the first set, if there already is a value in the parsing table for (nonterm,term), then we have a conflict and we print that to the screen. If not, and if the term is not epsilon, we add the rhs with the index of the production to the table at (nonterm,term). If that first set contained epsilon, then we need to also check the follow set. For each term in follow, if the term is not epsilon, we check whether there already is a value in the parsing table for (nonterm,term). If yes, we print that we have a conflict, if not, we add to the table. If that term is epsilon, then we add to (nonterm, dollarsign).

**Configuration:**

The configuration class contains the parsing table, list of production indexes, input stack and working stack. We have methods for getting the initial config, pushing, popping, checking if we should push, pop, or if the sequence is accepted. The initial config contains the input sequence in the input stack, and the working stack contains the start symbol. The push operation gets the value from the parsing table corresponding to (top\_of\_working\_stack, top\_of\_input\_stack), pops from the working stack, and if the value in the parse table does not contain epsilon, adds to the working stack the rhs from the table value, then adds the index of the production from the table value to the production indexes list

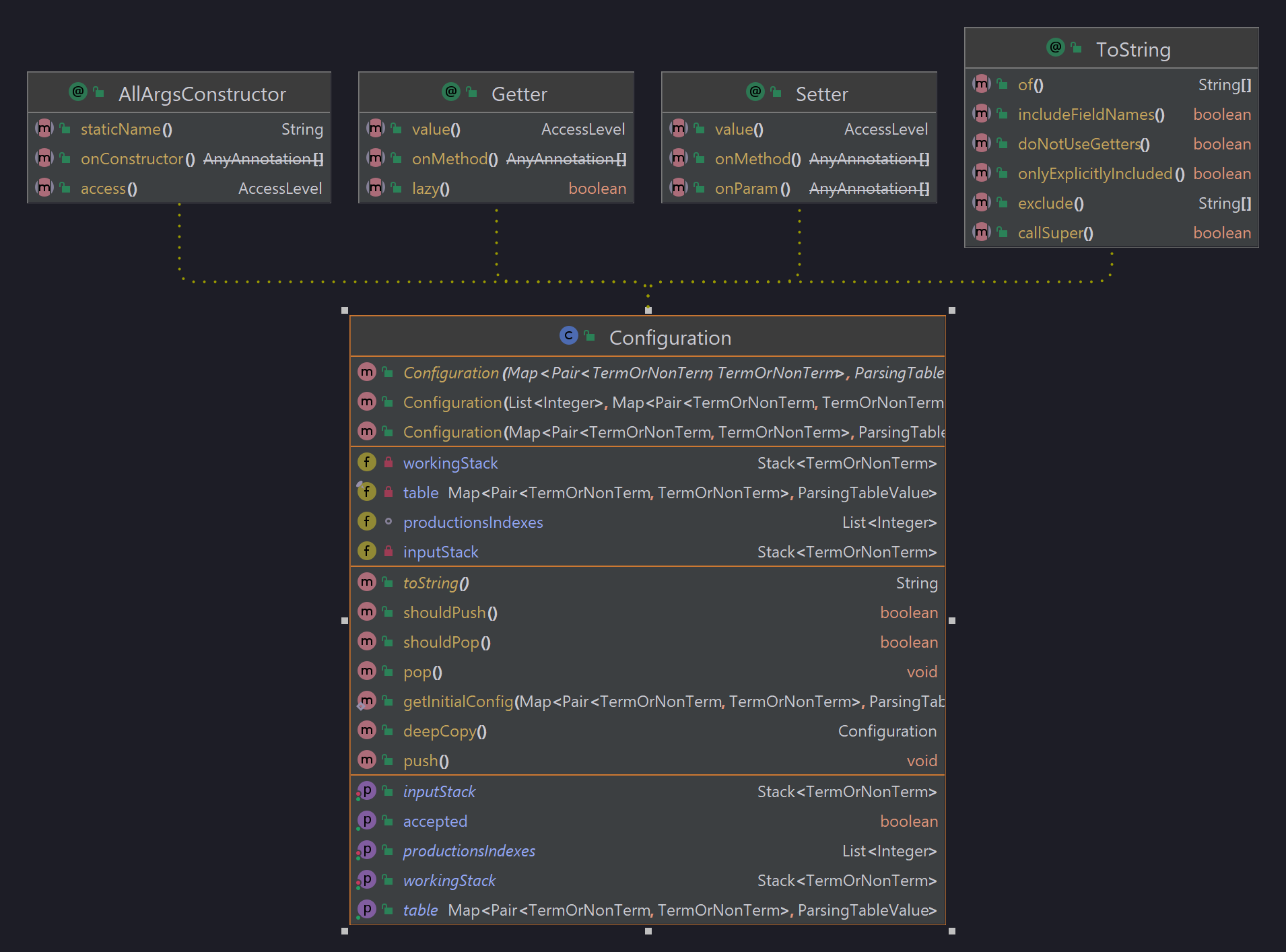
The pop operation just pops the input and working stack. The shouldPush method checks that the sequence is not accepted, that shouldPop is false and the value in the parsing table is not an error. The method shouldPop checks that the value in the parsing table corresponding to (top\_of\_working\_stack, top\_of\_input\_stack) is pop. In the same way, isAccepted checks that the value in the parsing table is Acc.

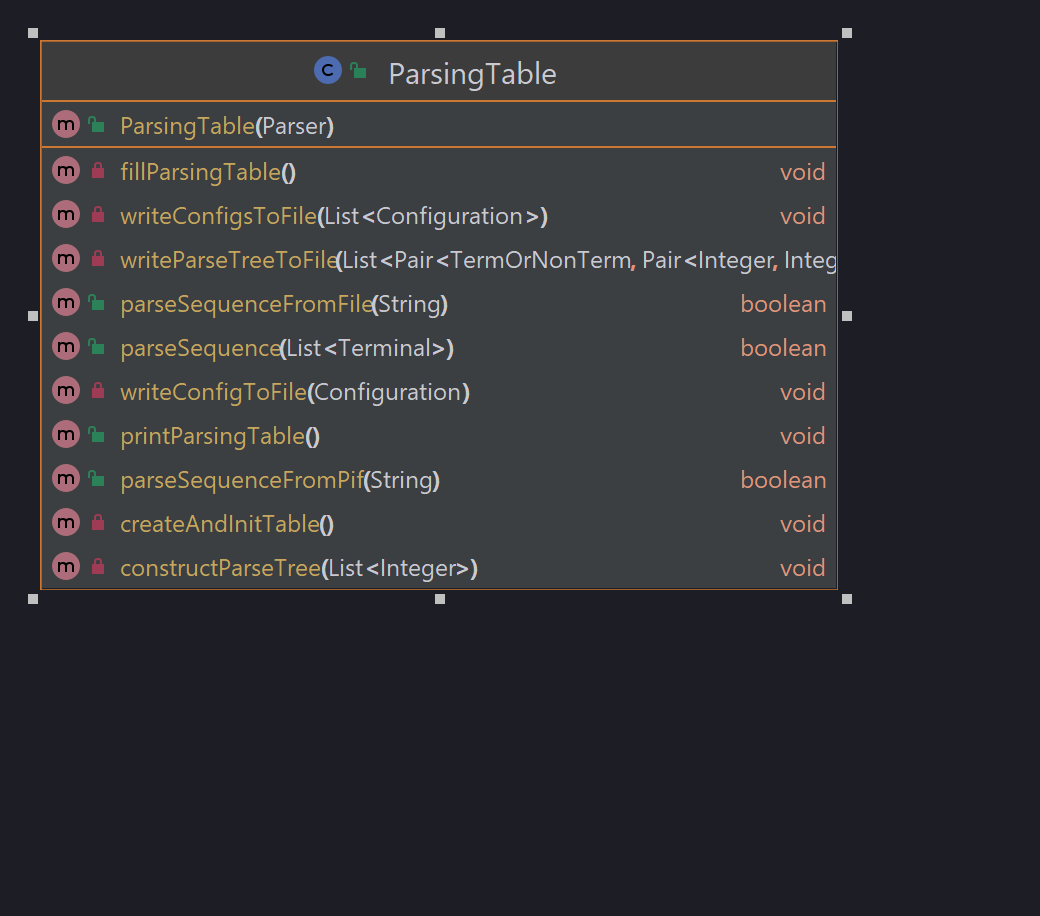
**Parsing sequence algorithm:**

We get the initial config. Set the Boolean flag to true, and while run, print the current configuration to a file. If should push, then push, else, if should pop, then pop, else if is not accepted, set the error flag to true and set the run flag to false, and if it is accepted, just set the run flag to false. Depending on the error flag, we print whether the sequence was accepted or not, and start building the parse tree.

**Building the parse tree – parent sibling representation:**

The data structure is a list of pairs of (TermOrNonTerm, Pair<Integer,Integer>), representing that for each entry in the list, we have the TermOrNonTerm, the index of the parent and the index of the left sibling. We use the value -1 for marking an invalid index, for the case of the root and the leftmost sibling which does not have a left sibling. We first add an entry for the starting symbol with parent -1 and left sibling -1. We also have a queue with indexes, and we first add index 0 corresponding to the root. Then we take each production corresponding to the production indexes found in the parsing sequence step, we get the current parent by popping from the queue, and for each element in the rhs of the production, we add to the table with that current parent, also making sure to put the left sibling as the index of the previous element. If the current element from the rhs is a nonterminal, we add to the queue, since those elements will be parent. Then we print the parse tree to the file.

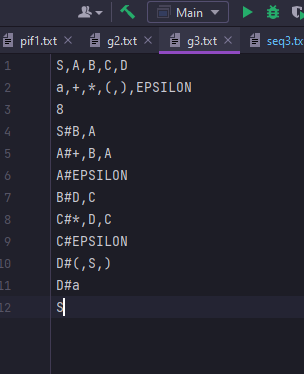


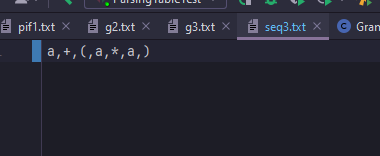


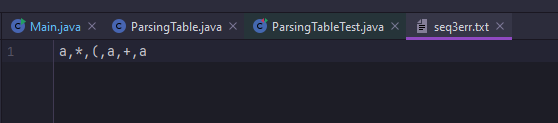
**Testing**: tests the parser accepts or not accepts different sequences from different grammars correctly

Input grammars and sequences:

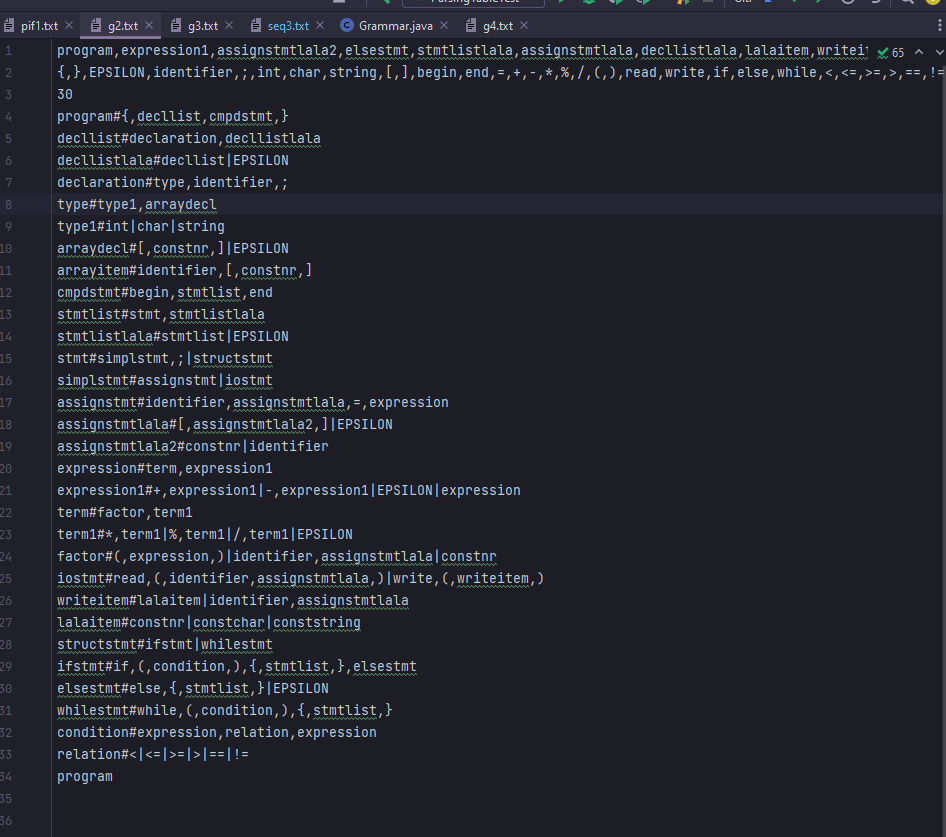
Grammar from seminary:



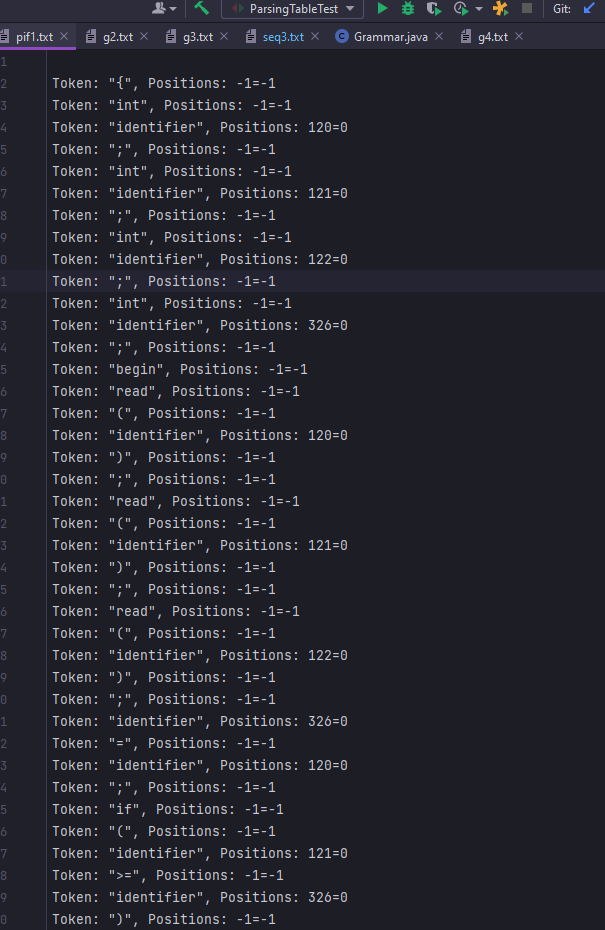




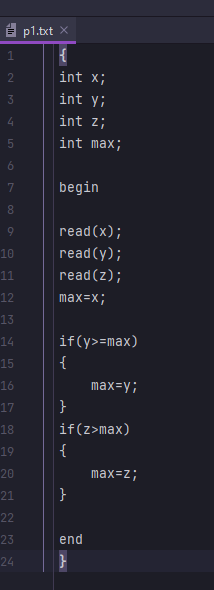
Language grammar:

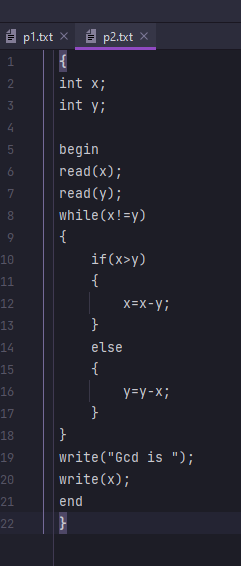
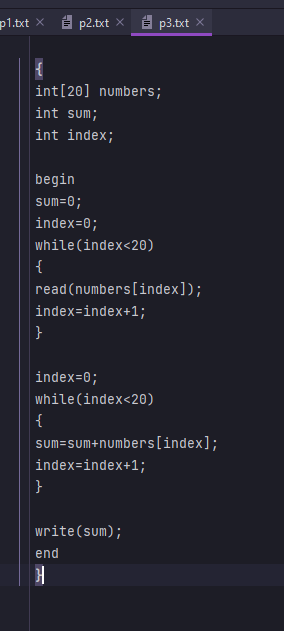
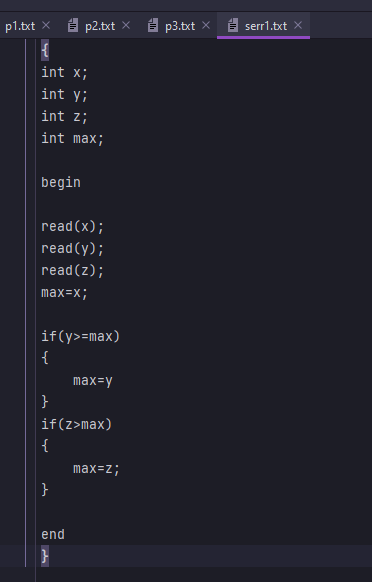
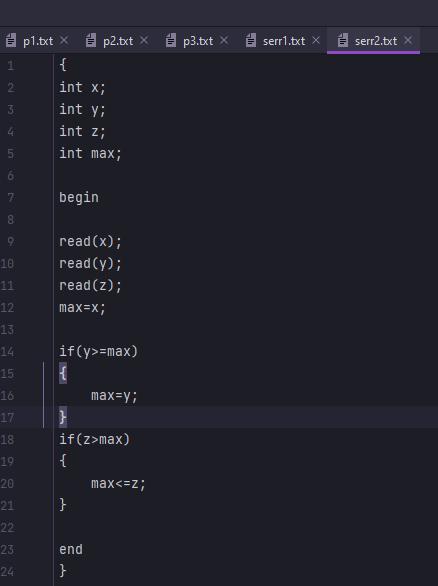


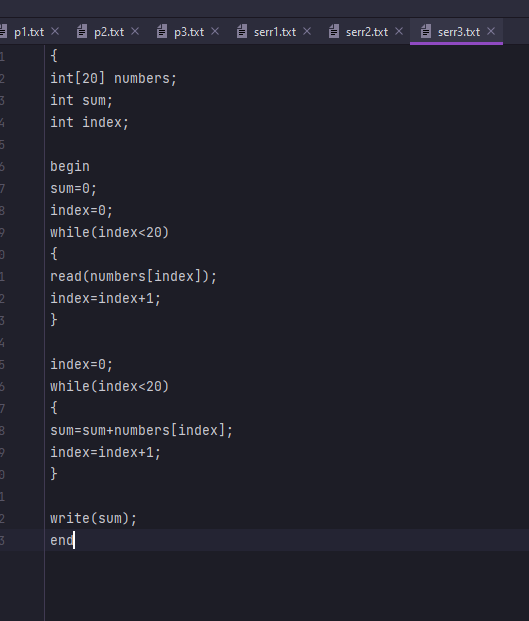
Pifs corresponding to programs from lab1a (p1.txt,p2.txt,p3.txt, and multiple types of syntax errors)



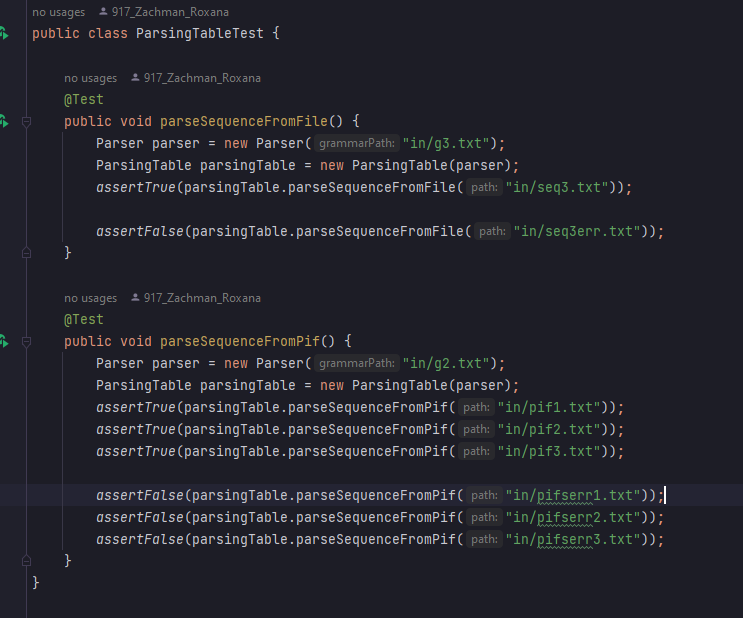
(I will put here the code for the programs, since it would be easier to check)





Tests:



For g3 and seq3 the config file and the parse tree file are:

